# DM545/DM871 <br> Linear and Integer Programming 

# MILP in Spreadsheets 

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## Outline

1. Example 1: Production Planning
2. Example 2: Diet Problem
3. Example: Budget Allocation
4. Example 4: Network Problem
5. Example 1: Production Planning
6. Example 2: Diet Problem
7. Example: Budget Allocation
8. Example 4: Network Problem

## Production Planning



Mathematical Model:

$$
\begin{aligned}
\max 16 x_{1}+10 x_{2} & \\
2 x_{1}+2 x_{2} & \leq 8 \\
2 x_{1}+x_{2} & \leq 6 \\
x_{1} & \geq 0 \\
x_{2} & \geq 0
\end{aligned}
$$

## General Model

$$
\begin{aligned}
& \max \sum_{j=1}^{n} c_{j} x_{j} \\
& \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, i=1, \ldots, m \\
& x_{j} \geq 0, j=1, \ldots, n
\end{aligned}
$$

$$
\begin{aligned}
\max c^{\top} x & \\
A x & \leq b \\
x & \geq 0
\end{aligned}
$$

$$
x \in \mathbb{R}^{n}, c \in \mathbb{R}^{n}, A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^{m}
$$

$$
\begin{aligned}
\max 16 x_{1}+10 x_{2} & \\
2 x_{1}+2 x_{2} & \leq 8 \\
2 x_{1}+x_{2} & \leq 6 \\
x_{1}, x_{2} & \geq 0
\end{aligned}
$$

$$
x_{1}, x_{2} \geq 0
$$

## Vector and Matrices in Excel

$$
\sum_{j=1}^{n} c_{j}=c_{1}+c_{2}+\ldots+c_{n}
$$

SUM(B5: B14)

## Scalar product

$$
\begin{aligned}
\mathbf{u} \cdot \mathbf{v} & =u_{1} v_{1}+u_{2} v_{2}+\ldots+u_{n} v_{n} \\
& =\sum_{j=1}^{n} u_{j} v_{j}
\end{aligned}
$$

SUMPRODUCT(B5: B14, C5: C14)

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## The Diet Problem (Blending Problems)

- Select a set of foods that will satisfy a set of daily nutritional requirement at minimum cost.
- Motivated in the 1930 s and 1940 s by US army.
- Formulated as a linear programming problem by George Stigler
- First linear programming problem
- (programming intended as planning not computer code)

min cost/weight
subject to nutrition requirements:
eat enough but not too much of Vitamin A eat enough but not too much of Sodium eat enough but not too much of Calories


## The Diet Problem

Suppose there are:

- 3 foods available, corn, milk, and bread,
- there are restrictions on the number of calories (between 2000 and 2250) and the amount of Vitamin A (between 5000 and 50,000 )

| Food | Corn | $2 \%$ Milk | Wheat bread |
| :--- | ---: | ---: | ---: |
| Vitamin A | 107 | 500 | 0 |
| Calories | 72 | 121 | 65 |
| Cost per serving | $\$ 0.18$ | $\$ 0.23$ | $\$ 0.05$ |

Rescaled

| Food | Corn | 2\% Milk | Wheat bread |
| :--- | ---: | ---: | ---: |
| Vitamin A | 983.33 | 2173.91 | 0 |
| Calories | 400 | 526.08 | 1300 |
| Cost per serving | $\$ 1$ | $\$ 1$ | $\$ 1$ |

$$
\begin{aligned}
\min \sum_{i \in F} c_{i} x_{i} & \\
\sum_{i \in F} a_{i j} x_{i} \geq N_{\operatorname{minj} j}, & \forall j \in N \\
\sum_{i \in F} a_{i j} x_{i} \leq N_{\max j}, & \forall j \in N \\
x_{i} \geq F_{\min i}, & \forall i \in F \\
x_{i} \leq F_{\max i}, & \forall i \in F
\end{aligned}
$$

## The History of Stigler's Diet Problem

- The linear program consisted of 9 equations in 77 variables
- Stigler, guessed an optimal solution using a heuristic method
- In 1947, the National Bureau of Standards used the newly developed simplex method to solve Stigler's model.
It took 9 clerks using hand-operated desk calculators 120 man days to solve for the optimal solution

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## Budget Allocation - Exercise 6 of Sheet 3

- A company has six different opportunities to invest money.
- Each opportunity requires a certain investment over a period of 6 years or less.

| Expected <br> Investment Cash <br> Flows and Net <br> Present Value |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | Opp. 1 | Opp. 2 | Opp. 3 | Opp. 4 | Opp. 5 | Opp. 6 |  |
| Year 1 | $-\$ 5.00$ | $-\$ 9.00$ | $-\$ 12.00$ | $-\$ 7.00$ | $-\$ 20.00$ | $-\$ 18.00$ |  |
| Year 2 | $-\$ 6.00$ | $-\$ 6.00$ | $-\$ 10.00$ | $-\$ 5.00$ | $\$ 6.00$ | $-\$ 15.00$ |  |
| Year 3 | $-\$ 16.00$ | $\$ 6.10$ | $-\$ 5.00$ | $-\$ 20.00$ | $\$ 6.00$ | $-\$ 10.00$ |  |
| Year 4 | $\$ 12.00$ | $\$ 4.00$ | $-\$ 5.00$ | $-\$ 10.00$ | $\$ 6.00$ | $-\$ 10.00$ |  |
| Year 5 | $\$ 14.00$ | $\$ 5.00$ | $\$ 25.00$ | $-\$ 15.00$ | $\$ 6.00$ | $\$ 35.00$ |  |
| Year 6 | $\$ 15.00$ | $\$ 5.00$ | $\$ 15.00$ | $\$ 75.00$ | $\$ 6.00$ | $\$ 35.00$ |  |
| NPV | $\$ 8.01$ | $\$ 2.20$ | $\$ 1.85$ | $\$ 7.51$ | $\$ 5.69$ | $\$ 5.93$ |  |

- The company has an investment budget that needs to be met for each year.
- It also has the wish of investing in those opportunities that maximize the combined Net Present Value (NPV) after the 6th year.


## Digression: What is the Net Present Value?

- $P$ : value of the original payment presently due
- the debtor wants to delay the payment for $t$ years,
- let $r$ be the market rate of return that the creditor would obtain from a similar investment asset
- the future value of $P$ is $F=P(1+r)^{t}$

Viceversa, consider the task of finding:

- the present value $P$ of $F=\$ 100$ that will be received in five years, or equivalently,
- which amount of money today will grow to $F=\$ 100$ in five years when subject to a constant discount rate? $P=\frac{F}{(1+r)^{t}}$
Assuming a $5 \%$ per year interest rate, it follows that

$$
P=\frac{F}{(1+r)^{t}}=\frac{\$ 100}{(1+0.05)^{5}}=\$ 78.35 .
$$

## Budget Allocation

Net Present Value calculation:
for each opportunity we calculate the NPV at time zero (the time of decision) as:

$$
P_{0}=\sum_{t=1}^{5} \frac{F_{t}}{(1+0.05)^{5}}
$$

| Expected <br> Investment Cash <br> Flows and Net <br> Present Value |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |

## Budget Allocation - Mathematical Model

- Let $B_{t}$ be the budget available for investments during the years $t=1 . .5$.
- Let $a_{t j}$ be the cash flow for opportunity $j$ and $c_{j}$ its NPV
- Task: choose a set of opportunities such that the budget is never exceeded and the expected return is maximized. Consider both the case of indivisible and divisible opportunities.

Variables $x_{j}=1$ if opportunity $j$ is selected and $x_{j}=0$ otherwise, $j=1 . .6$
Objective

$$
\max \sum_{j=1}^{6} c_{j} x_{j}
$$

Constraints

$$
\sum_{j=1}^{6} a_{t j} x_{j}+B_{t} \geq 0 \quad \forall t=1 . .5
$$

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## Min Cost Flow

A company produces the same product at two different factories, $A$ and $B$, and then the product must be shipped to two warehouses, $D$ and $E$, where either factory can supply either warehouse. A distribution center $C$ can be used to collect the product from $A$ and $B$ before shipping it to $E$. The distribution network has the following allowed connections: $\{A D, A B, A C, B C, C E, D E, E D\}$ and there are costs and bounds on the amount of product to ship through the connections.


## Minimum Cost Network Flows

Find cheapest flow through a network in order to satisfy demands at certain nodes from available supplier nodes.

## Variables:

$$
x_{i j} \in \mathbb{Z}_{0}^{+}
$$

## Objective:

$$
\min \sum_{i j \in A} c_{i j} x_{i j}
$$

Constraints: mass balance + flow bounds

$$
\begin{aligned}
& \sum_{j: i j \in A} x_{i j}-\sum_{j: j i \in A} x_{j i}=b(i) \quad \forall i \in V \\
& l_{i j} \leq x_{i j} \leq u_{i j}
\end{aligned}
$$



Example of mass balance constraints:

$$
x_{C E}-x_{A C}-x_{B C}=0
$$

## In Matrix Form

| i | $X_{e_{1}}$ $C_{e_{1}}$ | $X_{e_{2}}$ $C_{e_{2}}$ | . . . | $x_{i j}$ $c_{i j}$ |  | $X_{e_{m}}$ <br> $C_{e_{m}}$ |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | . | . . | - | . . | . | $=$ | $b_{1}$ |
| 2 1 | - | . | . . | . | $\ldots$ | . | $=$ | $b_{2}$ |
| : 1 | : | $\because$. |  |  |  |  | $=$ | : |
| $i$ | $-1$ | . | . . . | 1 | . . | . | $=$ | $b_{i}$ |
| : 1 | : | $\because$ |  |  |  |  | $=$ | : |
| $j$ | . | . | $\ldots$ | -1 | . . | . | $=$ | $b_{j}$ |
| $\vdots 1$ | . | $\bullet$. |  |  |  |  | $=$ | : |
| $n \quad 1$ |  |  |  |  |  |  | $=$ | $b_{n}$ |
| $e_{1}$ | 1 |  |  |  |  |  | $\leq$ | $u_{1}$ |
| $e_{2}$ |  | 1 |  |  |  |  | $\leq$ | $u_{2}$ |
| $\therefore 1$ | $\cdot$ | $\bullet$. |  |  |  |  | $\leq$ | : |
| $(i, j)$ |  |  |  | 1 |  |  | $\leq$ | $u_{i j}$ |
| : 1 | : | $\because$ |  |  |  |  | $\leq$ | $\vdots$ |
| $e_{m}$ |  |  |  |  |  | 1 | $\leq$ | $u_{m}$ |

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N node arc incidence matrix

