## DM545/DM871 - Linear and integer programming

Sheet 7, Spring 2024

Exercises with the symbol ${ }^{+}$are to be done at home before the class. Exercises with the symbol *will be tackled in class and should be at least read at home. The remaining exercises are left for self training after the exercise class.

## Exercise $1^{+}$

Consider the following three matrices:

$$
\left[\begin{array}{ccccc}
1 & 1 & -1 & 0 & 1 \\
1 & 0 & 0 & -1 & 0 \\
0 & 0 & 0 & 1 & 1 \\
0 & -1 & 1 & 0 & 0
\end{array}\right] \quad\left[\begin{array}{ccccc}
1 & 1 & 0 & 0 & 0 \\
-1 & 0 & 0 & 1 & -1 \\
0 & -1 & 1 & 0 & 1 \\
0 & 0 & 1 & 1 & 0
\end{array}\right] \quad\left[\begin{array}{ccc}
1 & 1 & 1 \\
1 & -1 & 1
\end{array}\right]
$$

For each of them say if it is totally unimodular and justify your answer.

## Exercise $2^{+}$

1. In class we stated that for the uncapaciteted facility location problem there are two formulations:

$$
\begin{gathered}
X=\left\{(x, y) \in \mathbb{R}_{+}^{m} \times \mathbb{B}^{1}: \sum_{i=1}^{m} x_{i} \leq m y, x_{i} \leq 1 \text { for } i=1, \ldots, m\right\} \\
P=\left\{(x, y) \in \mathbb{R}_{+}^{m} \times \mathbb{R}^{1}: x_{i} \leq y \text { for } i=1, \ldots, m, y \leq 1\right\}
\end{gathered}
$$

Prove that the polyhedron $P=\left\{\left(x_{1}, \ldots, x_{m}, y\right) \in \mathbb{R}^{m+1}: y \leq 1, x_{i} \leq y\right.$ for $\left.i=1, \ldots, m\right\}$ has integer vertices. [Hint: start by writing the constraint matrix and show that it is TUM.]
2. Consider the following (integer) linear programming problem:

$$
\begin{align*}
\min & c_{1} x_{1}+c_{2} x_{2}+c_{3} x_{3}+c_{4} x_{4} \\
& x_{3}+x_{4} \geq 10 \\
& x_{2}+x_{3}+x_{4} \geq 20 \\
& x_{1}+x_{2}+x_{3}+x_{4} \geq 30  \tag{1}\\
& x_{2}+x_{3} \geq 15 \\
& x_{1}, x_{2}, x_{3}, x_{4} \in \mathbb{Z}_{0}^{+}
\end{align*}
$$

The constraint matrix has consecutive 1's in each column. Matrices with consecutive 1's property for each column are totally unimodular. Show that this fact holds for the specific numerical example (1). That is, show first that the constraint matrix of the problem has consecutive 1 s in the columns and then that you can transform this matrix into one that you should recognize to be a TUM matrix. [Hint: rewrite the problem in standard form (that is, in equation form) and add a redundant row $\mathbf{0} \cdot \mathbf{x}=0$ to the set of constraints. Then perform elementary row operations to bring the matrix to a TUM form.]

## Exercise $3^{+}$

In class, we proved that the (mininum) vertex covering problem and the (maximum) matching problem are a weak dual pair. Prove that for bipartite graphs they, actually, are a strong dual pair.

## Exercise $4^{*}$ MILP Modeling

| Hour | Staffing requirement |
| :---: | :---: |
| 0 am to 6 am | 2 |
| 6 am to 8 am | 8 |
| 8 am to 11 am | 5 |
| 11 am to 2 pm | 7 |
| 2 pm to 4 pm | 3 |
| 4 pm to 6 pm | 4 |
| 6 pm to 8 pm | 6 |
| 8 pm to 10 pm | 3 |
| 10 pm to 12 pm | 1 |

Table 1:

Shift scheduling. The administrators of a department of a urban hospital have to organize the working shifts of nurses maintaining sufficient staffing to provide satisfactory levels of health care. Staffing requirements at the hospital during the whole day vary from hour to hour and are reported in Table 1. According to union agreements, nurses can work following one of the seven shift patterns in Table 2 each with its own cost.

| pattern | Hours of work | total hours | cost |
| :---: | :---: | :---: | ---: |
| 1 | 0 am to 6 am | 6 | 720 Dkk |
| 2 | 0 am to 8 am | 8 | 800 Dkk |
| 3 | 6 am to 2 pm | 8 | 740 Dkk |
| 4 | 8 am to 4 pm | 8 | 680 Dkk |
| 5 | 2 pm to 10 pm | 8 | 720 Dkk |
| 6 | 4 pm to 12 pm | 8 | 780 Dkk |
| 7 | 6 pm to 12 pm | 6 | 640 Dkk |

Table 2:

The department administrators would like to identify the assignment of nurses to working shifts that meets the staffing requirements and minimizes the total cost.

1. Model the problem as a MILP problem.
2. Use the results from Exesrcise 2 to show that the shift scheduling problem can be solved efficiently when formulated as a mathematical programming problem. (You do not need to find numerical results.)

## Exercise 5

Generalized Assignment Problem. Suppose there are $n$ types of trucks available to deliver products to $m$ clients. The cost of truck of type $i$ serving client $j$ is $c_{i j}$. The capacity of truck type $i$ is $Q_{i}$ and the demand of each client is $d_{j}$. There are $a_{i}$ trucks for each type. Formulate an IP model to decide how many trucks of each type are needed to satisfy all clients so that the total cost of doing the deliveries is minimized. If all the input data will be integer, will the solution to the linear programming relaxation always be integer?

## Exercise $6^{*}$ Network Flows: Problem of Representatives

A town has $r$ residents $R_{1}, R_{2}, \ldots, R_{r} ; q$ clubs $C_{1}, C_{2}, \ldots, C_{q}$; and $p$ political parties $P_{1}, P_{2}, \ldots, P_{p}$. Each resident is a member of at least one club and belongs to exactly one political party. Each club must nominate one of its members to represent it on the town's governing council so that the number of council members belonging to the political party $P_{k}$ is at most $u_{k}$. Is it possible to find a council that satisfies this "balancing" property?
Show how to formulate this problem as a maximum flow problem.

